

AN L -SYSTEM ON THE SMALL WITT DESIGN

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ABSTRACT. We construct a 12-uniform hypergraph on n vertices with size $(n/12)^6$ which satisfies $|F \cap F'| \in \{0, 1, 2, 3, 4, 6\}$ for all distinct edges F and F' .

1. INTRODUCTION

Let k and n be positive integers and let $L \subset \{0, 1, \dots, k-1\}$. A k -uniform hypergraph \mathcal{F} is called a (k, L) -system (or an L -system for short) if $|F \cap F'| \in L$ holds for all distinct $F, F' \in \mathcal{F}$. Let $m(n, k, L)$ be the maximum size of (k, L) -systems on n vertices. If there exist positive constants α, c, c' , and n_0 depending only on k and L such that $cn^\alpha < m(n, k, L) < c'n^\alpha$ holds for $n > n_0$, then we define $\alpha(k, L) = \alpha$ and we say that (k, L) -systems have exponent α . In [4] the following upper bound for the size of (k, L) -systems is obtained.

Theorem 1. *For $n > n_0(k, L)$ it follows that*

$$m(n, k, L) \leq \prod_{l \in L} \frac{n-l}{k-l}.$$

In particular, the above upper bound gives $\alpha(k, L) \leq |L|$ if $\alpha(k, L)$ exists. In this note, we construct some (k, L) -systems satisfying $\alpha(k, L) = |L|$. Among other results, we show that $\alpha(12, \{0, 1, 2, 3, 4, 6\}) = 6$. The corresponding system is related to the small Witt design $S(5, 6, 12)$ and our construction uses the embedding of the design into $PG(5, 3)$.

In the next section we explain our main idea by solving a toy problem. Then we state the general construction scheme in section 3. We deal with the L -systems related to the small Witt designs in section 4, and some other L -systems with geometric structures in section 5. Finally in section 6 we consider “intersection structures” which control L -systems. The results in section 6 together with previously known results give the complete tables of exponents of (k, L) -systems for $k = 11$ and $k = 12$. These tables are presented in the Appendix and settle all the open cases in [7].

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2. A TOY PROBLEM

As a toy problem, let us consider a $(7, \{0, 1, 3\})$ -system. Let A be the following 3×7 matrix over $GF(2)$:

$$A = \begin{pmatrix} 1 & 0 & 0 & 1 & 1 & 0 & 1 \\ 0 & 1 & 0 & 1 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 & 1 & 1 & 1 \end{pmatrix}.$$

This is a parity check matrix of the Hamming $[7, 4, 3]_2$ -code. This matrix has the following properties.

- (i) Any two columns are linearly independent over $GF(2)$.
- (ii) For any two columns c_p, c_q of A , the subspace spanned by c_p, c_q contains precisely three columns c_p, c_q and c_r of A . In fact, $c_r = c_p + c_q$.

Let us choose 3 columns of A . There are $\binom{7}{3} = 35$ choices. Among those choices, $\binom{7}{2} / \binom{3}{2} = 7$ of them span 2-dimensional subspaces and the other 28 span the entire 3-dimensional space. The triples (the indices of 3 columns) corresponding to 2-dimensional subspaces form the Steiner triple system $S(2, 3, 7) = \{124, 135, 167, 236, 257, 347, 456\}$ or the Fano plane.

Now we shall construct a $(7, \{0, 1, 3\})$ -system \mathcal{F} using the matrix A . This is going to be a 7-partite hypergraph on the vertex partition $X = V_1 \cup \dots \cup V_7$ where $V_i \cong \mathbb{F}_2^d$ for $1 \leq i \leq 7$. For $a, b, c \in \mathbb{F}_2^d$, define the ordered 7-tuple $F(a, b, c)$ by

$$F(a, b, c) = (a, b, c)A = (a, b, c, a + b, a + c, b + c, a + b + c) \in (\mathbb{F}_2^d)^7.$$

Then define

$$\mathcal{F} = \{F(a, b, c) : a, b, c \in \mathbb{F}_2^d\}.$$

Let us check that \mathcal{F} is a $\{0, 1, 3\}$ -system. Choose $F = F(a, b, c), F' = F(a', b', c') \in \mathcal{F}$. Suppose that $i := |F \cap F'| \geq 2$ and $F = F'$ on $V_p \cup V_q$. Then by (ii), $F = F'$ on V_r as well, where r is given by $c_r = c_p + c_q$. This means $i \geq 3$ (and $i = 2$ cannot happen). Next suppose that $i \geq 4$. Then by (ii) we can choose p, q, s where $1 \leq p < q < s \leq 7$ such that $F = F'$ on $V_p \cup V_q \cup V_s$ and the 3×3 minor matrix $B = (c_p, c_q, c_s)$ of A is non-singular. Since $(a - a', b - b', c - c')B = 0$ we have $(a, b, c) = (a', b', c')$ and $F = F'$. This means $i = 7$ (and $i = 4, 5$, or 6 cannot occur). Therefore, \mathcal{F} is a $(7, \{0, 1, 3\})$ -system.

The size of this hypergraph is $|\mathcal{F}| = (n/7)^3$ where $n = |V(\mathcal{F})| = 7 \cdot 2^d$. We will generalize the construction in the next section.

In this particular case, we can construct a larger $(7, \{0, 1, 3\})$ -system by taking \mathcal{F} to be the set of projective planes in $PG(d, 2)$. Then,

$$|\mathcal{F}| = \frac{(2^{d+1} - 1)(2^{d+1} - 2)(2^{d+1} - 4)}{(2^3 - 1)(2^3 - 2)(2^3 - 4)} = \frac{n(n-1)(n-3)}{7 \cdot 6 \cdot 4}$$

where $n = |PG(d, 2)| = 2^{d+1} - 1$. This size attains the upper bound in Theorem 1, so this construction is best possible.

3. GENERATING MATRIX FOR A (k, L) -SYSTEM

A $(t, b, k)_q$ -matrix is a $(t+1) \times k$ matrix over $GF(q)$ satisfying the following properties:

- (P1) Any t columns are linearly independent over $GF(q)$.
- (P2) For any t columns, the t -dimensional subspace spanned by these columns contains precisely b columns of A .

The matrix in the previous section was a $(2, 3, 7)_2$ -matrix. For a $(t, b, k)_q$ -matrix A , there are $\binom{k}{b}$ ways of taking a $(t+1) \times b$ minor matrix of A . Among $\binom{k}{b}$ ways, $\binom{k}{t}/\binom{b}{t}$ of them have rank t , and the others have rank $t+1$. Each of those rank t minor matrices gives a b -set (block) consisting of the indices of corresponding columns. The set of these blocks is a Steiner system $S(t, b, k)$. To represent this situation, we say that a $(t, b, k)_q$ -matrix supports $S(t, b, k)$.

Theorem 2. *If there exists a $(t, b, k)_q$ -matrix then there exists a (k, L) -system on n -vertices with size $(n/k)^{t+1}$ where $L = \{0, 1, \dots, t-1, b\}$.*

Proof. Let $A = (a_{ij}) = (c_1, \dots, c_k)$ ($1 \leq i \leq t+1$, $1 \leq j \leq k$) be a $(t, b, k)_q$ -matrix. We shall construct a (k, L) -system \mathcal{F} which is k -partite on the vertex partition $X = V_1 \cup \dots \cup V_k$ where $V_i \cong \mathbb{F}_q^d$ for $1 \leq i \leq k$. For $(x_1, \dots, x_{t+1}) \in (\mathbb{F}_q^d)^{t+1}$, let us define the k -set $F(x_1, \dots, x_{t+1}) \in (\mathbb{F}_q^d)^k$ by setting

$$F(x_1, \dots, x_{t+1}) = (x_1, \dots, x_{t+1})A = \left(\sum_{i=1}^{t+1} a_{ij}x_i \right)_{j=1}^k \in V_1 \times \dots \times V_k.$$

Then define

$$(1) \quad \mathcal{F} = \{F(x_1, \dots, x_{t+1}) : x_1, \dots, x_{t+1} \in \mathbb{F}_q^d\}.$$

By construction, \mathcal{F} is k -partite and k -uniform. Let us check that \mathcal{F} is an L -system. Choose two edges $F = F(x_1, \dots, x_{t+1})$ and $F' = F(x'_1, \dots, x'_{t+1})$ of \mathcal{F} . Let $i := |F \cap F'|$ and let I be the corresponding i -set such that $F = F'$ on $\bigcup_{i \in I} V_i$. Set $y = (x_1 - x'_1, \dots, x_{t+1} - x'_{t+1})$ then $y \cdot c_i = 0$ holds for $i \in I$.

Suppose that $b \geq i \geq t$. Then by (P2) there exists $B \supset I$, $|B| = b$ such that $c_j = \sum_{i \in I} \gamma_{ij} c_i$ holds for $j \in B$ where $\gamma_{ij} \in \mathbb{F}_q$. Thus for $j \in B$ it follows that $y \cdot c_j = y \cdot \sum_{i \in I} \gamma_{ij} c_i = \sum_{i \in I} \gamma_{ij} (y \cdot c_i) = 0$. So we have $F = F'$ on V_i , $i \in B$. This means $i = b$.

Next suppose that $i \geq b + 1$. Then by (P2) we can choose $t + 1$ columns from $\{c_i : i \in I\}$ so that the corresponding $(t + 1) \times (t + 1)$ minor matrix C is non-singular. Then we have $yC = 0$, which implies $y = 0$, i.e., $F = F'$. This means $i = k$, which concludes that \mathcal{F} is an L -system.

The (k, L) -system \mathcal{F} has $n = k \cdot q^d$ vertices and size $|\mathcal{F}| = (q^d)^{t+1} = (n/k)^{t+1}$. \square

It is now appropriate to say that a $(t, b, k)_q$ -matrix is a generating matrix for a (k, L) -system where $L = \{0, 1, \dots, t - 1, b\}$. The row vectors of a $(t, b, k)_q$ -matrix span a $(t + 1)$ -dimensional subspace in a k -dimensional space. This fact together with (P2) implies that a $(t, b, k)_q$ -matrix is a parity check matrix of a $[k, k - t - 1, t + 1]_q$ -code.

Note that Theorem 1 and Theorem 2 imply that $\alpha(k, L) = |L|$ where $L = \{0, 1, \dots, t - 1, b\}$ if a $(t, b, k)_q$ -matrix exists.

4. A (12, 012346)-SYSTEM AND ITS DERIVED SYSTEM

Let us construct a $(5, 6, 12)_3$ -matrix A . We use a geometric structure due to Havlicek[10] originated by Coxeter[2]. Let $\varphi : PG(2, 3) \rightarrow PG(5, 3)$ be the Veronese mapping, that is, $\varphi(x, y, z) = (x^2, xy, xz, y^2, yz, z^2)$. Choose a line ℓ in $PG(2, 3)$, say, $\ell = \{001, 010, 011, 012\}$. Its Veronese image $\Gamma := \varphi(\ell) = \{000001, 000100, 000111, 000121\}$ is a planar quadrangle with the diagonal triangle $\Delta = \{000101, 000211, 000221\}$. Now we take $(\varphi(PG(2, 3)) - \Gamma) \cup \Delta$ as the 12 column vectors of A . More concretely, we have

$$A = \begin{pmatrix} 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 2 & 0 & 1 & 2 & 2 & 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 2 & 2 & 1 & 2 & 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 1 & 0 & 1 & 1 & 1 & 1 & 1 & 2 & 2 \\ 0 & 0 & 0 & 0 & 0 & 2 & 2 & 1 & 1 & 0 & 1 & 2 \\ 0 & 0 & 1 & 0 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \end{pmatrix}.$$

The 12 points in $PG(5, 3)$ determined by A have many interesting geometric properties. In particular, A is a $(5, 6, 12)_3$ -matrix. See [2, 10] for details.

The matrix is a parity check matrix of a $[12, 6, 6]_3$ -code. Noting that this is the unique extended ternary Golay code G_{12} , the following

standard parity check matrix B of G_{12} is also a $(5, 6, 12)_3$ -matrix.

$$B = \begin{pmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & 1 \\ 0 & 1 & 0 & 0 & 0 & 0 & 2 & 0 & 1 & 2 & 2 & 1 \\ 0 & 0 & 1 & 0 & 0 & 0 & 2 & 1 & 0 & 1 & 2 & 2 \\ 0 & 0 & 0 & 1 & 0 & 0 & 2 & 2 & 1 & 0 & 1 & 2 \\ 0 & 0 & 0 & 0 & 1 & 0 & 2 & 2 & 2 & 1 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 1 & 2 & 1 & 2 & 2 & 1 & 0 \end{pmatrix}.$$

These matrices support the Witt design $S(5, 6, 12)$.

By deleting the top row and the left-most column from A or B , we obtain a $(4, 5, 11)_3$ -matrix. This corresponds to a $(11, \{0, 1, 2, 3, 5\})$ -system, $S(4, 5, 11)$, and the perfect ternary Golay code G_{11} .

Consequently, we have the following bounds for the size of L -systems on the small Witt designs. (Lower bounds follow from the constructions above and Theorem 2, while upper bounds follow from Theorem 1.)

Theorem 3. *For n sufficiently large, we have*

$$\left(\frac{n}{12}\right)^6 \leq m(n, 12, \{0, 1, 2, 3, 4, 6\}) \leq \frac{n^6}{12 \cdot 11 \cdot 10 \cdot 9 \cdot 8 \cdot 6},$$

$$\left(\frac{n}{11}\right)^5 \leq m(n, 11, \{0, 1, 2, 3, 5\}) \leq \frac{n^5}{11 \cdot 10 \cdot 9 \cdot 8 \cdot 6}.$$

By deleting the first two rows and the first two columns from B , we obtain a $(3, 4, 10)_3$ -matrix C which supports $S(3, 4, 10)$. This gives $(n/10)^4$ as lower bound for a $(10, \{0, 1, 2, 4\})$ -system on n vertices. In [5], $n^4/65610$ is obtained as lower bound by a construction using sections of elliptic quadrics over $GF(3)$. We can use quadrics to construct generating matrices. In fact, each column vector (x, y, z, w) of the matrix C satisfies $xy - xz + xw - yz - yw + zw = 0$, which means the 10 points are on the Möbius plane. We will extend the construction in the next section.

5. MORE EXAMPLES OF A GENERATING MATRIX

Let q be a prime power. Here we present some examples of a $(t, b, k)_q$ -matrix and its corresponding (k, L) -system on n vertices. Similar constructions are also given in [5], but our constructions are simpler and give better lower bounds for (k, L) -systems.

5.1. Affine plane. Let $d \geq m$ and set $n = q^d$, $k = q^m$, $L = \{0, 1, q, \dots, q^{m-1}\}$. Let \mathcal{F} be the set of m -dimensional affine subspaces in \mathbb{F}_q^d . This is a (k, L) -system on n vertices with size

$$|\mathcal{F}| = \frac{q^d(q^d - 1)(q^d - q) \cdots (q^d - q^{m-1})}{q^m(q^m - 1)(q^m - q) \cdots (q^m - q^{m-1})} = \prod_{\ell \in L} \frac{n - \ell}{k - \ell}.$$

Comparing to Theorem 1, it follows that $|\mathcal{F}| = m(n, k, L)$, in particular, $\alpha(q^m, \{0, 1, q, \dots, q^{m-1}\}) = m$.

Now let $m = 2$. Then the above \mathcal{F} is a $(q^2, \{0, 1, q\})$ -system. By taking q^2 points of an affine plane as column vectors, we obtain a $(2, q, q^2)_q$ -matrix. For example, the case $q = 3$ gives the following $(2, 3, 9)_3$ -matrix:

$$\begin{pmatrix} 0 & 1 & 2 & 0 & 1 & 2 & 0 & 1 & 2 \\ 0 & 0 & 0 & 1 & 1 & 1 & 2 & 2 & 2 \\ 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \end{pmatrix}.$$

Taking $m = 3$, $q = 2$, we get a $(3, 4, 8)_2$ -matrix:

$$\begin{pmatrix} 0 & 1 & 0 & 1 & 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 1 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \end{pmatrix}.$$

5.2. Projective plane. Let $d \geq m$ and set $n = q^{d+1} - 1$, $k = \frac{q^{m+1}-1}{q-1}$, $L = \{0, 1, q+1, \dots, \frac{q^m-1}{q-1}\}$. Let \mathcal{F} be the set of m -dimensional projective subspaces in $PG(d, q)$. This is a (k, L) -system on n vertices with size

$$|\mathcal{F}| = \frac{(q^{d+1} - 1)(q^{d+1} - q) \cdots (q^{d+1} - q^m)}{(q^{m+1} - 1)(q^{m+1} - q) \cdots (q^{m+1} - q^m)},$$

which gives $|\mathcal{F}| = \Omega(n^{m+1})$ as d (and hence n) grows. This together with Theorem 1 implies $\alpha(\frac{q^{m+1}-1}{q-1}, \{0, 1, q+1, \dots, \frac{q^m-1}{q-1}\}) = m + 1$.

Now let $m = 2$. Then the above construction gives a $(q^2 + q + 1, \{0, 1, q+1\})$ -system. By taking $p^2 + p + 1$ points of the projective plane as column vectors, we get a $(2, q+1, q^2 + q + 1)_q$ -matrix. The $(2, 3, 7)_2$ -matrix in section 2 is one of the examples.

5.3. Möbius plane. Set $k = q^2 + 1$, $L = \{0, 1, 2, q+1\}$, and let

$$\mathcal{M} = \{(x, y, z, w) \in PG(3, q) : f(x, y) + zw = 0\},$$

where $f(x, y)$ is an \mathbb{F}_q -irreducible quadratic form. Then $|\mathcal{M}| = q^2 + 1$ and we obtain a $(3, q+1, q^2 + 1)_q$ -matrix where the column vectors

come from \mathcal{M} (cf. Example 26.5 in [11]). Therefore we have $\alpha(q^2 + 1, \{0, 1, 2, q + 1\}) = 4$.

For example, setting $q = 3$, we consider $x^2 + y^2 + zw = 0$ in $PG(3, 3)$, which has 10 solutions giving a $(3, 4, 10)_3$ -matrix:

$$\begin{pmatrix} 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 & 0 & 0 & 1 & 1 & 2 & 2 \\ 0 & 1 & 1 & 2 & 1 & 2 & 1 & 2 & 1 & 2 \\ 1 & 0 & 2 & 1 & 2 & 1 & 1 & 2 & 1 & 2 \end{pmatrix}.$$

The following examples are a $(3, 5, 17)_4$ -matrix ($q = 4$ and $\mathbb{F}_4 = \mathbb{F}_2(\beta)$ where $\beta^2 + \beta + 1 = 0$, and $f(x, y) = x^2 + \beta xy + y^2$):

$$\begin{pmatrix} 0 & 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 & 1 & 0 & 0 & 0 & 1 & 1 & 1 & \beta & \beta & \beta & \beta^2 & \beta^2 & \beta^2 \\ 0 & 1 & 1 & \beta & \beta^2 & 1 & \beta & \beta^2 & 1 & \beta & \beta^2 & 1 & \beta & \beta^2 & 1 & \beta & \beta^2 \\ 1 & 0 & 1 & \beta^2 & \beta & 1 & \beta^2 & \beta & \beta & 1 & \beta^2 & 1 & \beta^2 & \beta & \beta & 1 & \beta^2 \end{pmatrix},$$

and a $(3, 6, 26)_5$ -matrix ($q = 5$ and $f(x, y) = x^2 + xy + y^2$):

$$\begin{pmatrix} 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & 2 & 2 & 2 & 2 & 3 & 3 & 3 & 3 & 4 & 4 & 4 & 4 \\ 0 & 1 & 1 & 2 & 3 & 4 & 1 & 2 & 3 & 4 & 1 & 2 & 3 & 4 & 1 & 2 & 3 & 4 & 1 & 2 & 3 & 4 & 1 & 2 & 3 & 4 \\ 1 & 0 & 4 & 2 & 3 & 1 & 4 & 2 & 3 & 1 & 2 & 1 & 4 & 3 & 3 & 4 & 1 & 2 & 2 & 1 & 4 & 3 & 4 & 2 & 3 & 1 \end{pmatrix}.$$

Let \mathcal{F} be a (k, L) -system on n vertices constructed by a $(3, q + 1, q^2 + 1)_q$ -matrix. Then we have $|\mathcal{F}| = \binom{n}{k}^4$ where $n = kq^d$, $k = q^2 + 1$. For fixed q we have $n \rightarrow \infty$ as d grows, and

$$\lim_{n \rightarrow \infty} \frac{|\mathcal{F}|}{\prod_{\ell \in L} \frac{n-\ell}{k-\ell}} = \frac{k(k-1)(k-2)(k-q-1)}{k^4} = \frac{q^2(q^2-1)(q^2-q)}{(q^2+1)^3} =: C_q.$$

Moreover we have $\lim_{q \rightarrow \infty} C_q = 1$. Therefore our construction is asymptotically best possible in this sense. In [5] an (n, k, L) -system ($k = q^2 + 1$, $L = \{0, 1, 2, q + 1\}$) \mathcal{F}' with $\lim_{n \rightarrow \infty} |\mathcal{F}'| / \prod_{\ell \in L} \frac{n-\ell}{k-\ell} = C'_q$ is constructed. In [5] and [3] they claim that $\lim_{q \rightarrow \infty} C'_q = 1$, but it seems that their construction only gives $\lim_{q \rightarrow \infty} C'_q = \frac{1}{2}$.

5.4. Unital. Set $k = q^3 + 1$, $L = \{0, 1, q + 1\}$, and let

$$\mathcal{U} = \{(x, y, z) \in PG(2, q^2) : x^{q+1} + y^{q+1} + z^{q+1} = 0\}.$$

Then $|\mathcal{U}| = q^3 + 1$ and we obtain a $(2, q + 1, q^3 + 1)_{q^2}$ -matrix (cf. Example 26.8 in [11], [1]). Therefore we have $\alpha(q^3 + 1, \{0, 1, q + 1\}) = 3$.

We say that $\mathcal{I} \subset 2^{[k]}$ is an intersection structure of a (k, L) -system if \mathcal{I} is a closed L -system whose rank is $\text{rank}(k, L)$. A generator set \mathcal{I}^* of \mathcal{I} is the collection of all maximal elements of \mathcal{I} , that is

$$\mathcal{I}^* := \{I \in \mathcal{I} : \nexists I' \in \mathcal{I} \text{ such that } I \subset I', I \neq I'\}.$$

We can retrieve \mathcal{I} from \mathcal{I}^* by taking all possible intersections.

For a family $\mathcal{F} \subset \binom{[n]}{k}$ and an edge $F \in \mathcal{F}$ define

$$\mathcal{I}(F, \mathcal{F}) := \{F \cap F' : F' \in \mathcal{F} - \{F\}\} \subset 2^F.$$

Moreover, if \mathcal{F} is k -partite with k -partition $[n] = X_1 \cup \dots \cup X_k$ then we define the projection $\pi(I)$ of $I \in \mathcal{I}(F, \mathcal{F})$ by $\pi(I) := \{i : I \cap X_i \neq \emptyset\} \subset 2^{[k]}$ and set $\pi(\mathcal{I}(F, \mathcal{F})) := \{\pi(I) : I \in \mathcal{I}(F, \mathcal{F})\}$. Füredi[8] proved the following fundamental result.

Theorem 4. *Given $k \geq 2$ and $L \subset \{0, 1, \dots, k-1\}$ there exists a positive constant $c = c(k, L)$ such that every (k, L) -system $\mathcal{F} \subset \binom{[n]}{k}$ contains a k -partite subfamily $\mathcal{F}^* \subset \mathcal{F}$ with k -partition $[n] = X_1 \cup \dots \cup X_k$ satisfying (1)–(3).*

- (1) $|\mathcal{F}^*| > c|\mathcal{F}|$.
- (2) For any two edges $F_1, F_2 \in \mathcal{F}^*$, $\pi(\mathcal{I}(F_1, \mathcal{F}^*)) = \pi(\mathcal{I}(F_2, \mathcal{F}^*))$.
- (3) For all $F \in \mathcal{F}^*$, $\mathcal{I}(F, \mathcal{F}^*)$ is a closed L -system.

In the above situation, we say that $\mathcal{I}(F, \mathcal{F}^*)$ is the intersection structure of \mathcal{F}^* . Let us see how the rank is related to the exponent of a (k, L) -system. Set $\mathcal{I} = \pi(\mathcal{I}(F, \mathcal{F}^*))$ and $t = \text{rank}(\mathcal{I})$, and consider \mathcal{F}^* in the above theorem. We can find some $A \in \binom{[k]}{t}$ such that $A \notin \Delta_t(\mathcal{I})$. Then for every $B \in \prod_{a \in A} V_a$ with $\pi(B) = A$ there is at most one member F of the family \mathcal{F}^* such that $B \subset F$. Thus the size $|\mathcal{F}^*|$ is at most the number of choices for B , that is $\prod_{a \in A} |V_a| = O(n^t)$. In other words, if $\alpha(k, L)$ exists then we have

$$\alpha(k, L) \leq \text{rank}(k, L).$$

On the other hand, Füredi[9] conjectures that

$$\alpha(k, L) > \text{rank}(k, L) - 1.$$

This is true if $\text{rank}(k, L) = 2$ (cf. [9]). As we will see in the next subsections, the conjecture is also true if $k \leq 12$ for all L . If there exists a Steiner system $S(t, b, k)$ then we have $\text{rank}(k, L) = t + 1$ for $L = [0, t-1] \cup \{b\}$. Rödl and Tengan[12] found a construction which shows $\alpha(k, L) > t$ in this situation. However there is no general lower bound known for $\alpha(k, L)$ in terms of $\text{rank}(k, L)$.

6.1. **The case $k = 11$.** Let $L_0 = \{0, 1, 2, 3, 5\}$. We consider (k, L) -systems with $k = 11$ and L listed below (16 cases).

- (I) $L_0, L_0 \cup \{6\}, L_0 \cup \{8\}, L_0 \cup \{9\}, L_0 \cup \{6, 8\}, L_0 \cup \{6, 9\}, L_0 \cup \{8, 9\},$
 $L_0 \cup \{6, 8, 9\}, L_0 \cup \{11\}, L_0 \cup \{6, 11\}, L_0 \cup \{8, 11\}, L_0 \cup \{9, 11\},$
 $L_0 \cup \{6, 8, 11\}, L_0 \cup \{6, 9, 11\}, L_0 \cup \{8, 9, 11\}, L_0 \cup \{6, 8, 9, 11\}.$

By computer search, we found that $\text{rank}(11, L) = 5$ for all L in (I), and the Steiner system $S(4, 5, 11)$ is the unique generator set of the corresponding intersection structure. As we saw in section 4 that $\alpha(11, L_0) = 5$, we now have $\alpha(11, L) = 5$ for all L in (I).

Next we consider the case $L = \{0, 1, 2, 3, 5, 7\}$. By computer search, we found that $\text{rank}(11, L) = 5$ and there are precisely two intersection structures — one is generated by $S(4, 5, 11)$ and the other is \mathcal{I}_{11} described below.

Let $J_i = \{2i, 2i + 1\}$ for $i = 1, \dots, 5$ and set

$$\mathcal{P} = \{\{1\} \cup J_a \cup J_b \cup J_c : \{a, b, c\} \in \binom{[5]}{3}\} \subset \binom{[11]}{7},$$

$$\mathcal{Q} = \{\{j_1, j_2, j_3, j_4, j_5\} : j_i \in J_i \text{ and } \sum j_i = \text{even}\} \subset \binom{[11]}{5}.$$

Then we define $\mathcal{I}_{11}^* = \mathcal{P} \cup \mathcal{Q}$ which is the generator set of \mathcal{I}_{11} .

Now we construct a 11-partite $(11, L)$ -system \mathcal{F} whose intersection structure is \mathcal{I}_{11} . Let A be the following generating matrix over $GF(2)$:

$$A = \begin{pmatrix} 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 \\ 0 & 0 & 0 & 0 & 0 & 1 & 1 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\ 0 & 1 & 1 & 0 & 0 & 1 & 1 & 1 & 1 & 1 & 1 \\ 1 & 0 & 1 & 0 & 1 & 0 & 1 & 0 & 1 & 0 & 1 \end{pmatrix}.$$

This is not a $(t, b, k)_q$ -matrix, but we can define the family \mathcal{F} on $V_1 \cup \dots \cup V_{11}$ where $V_i \cong \mathbb{F}_2^d$, by (1) in section 3.

Let c_i be the i -th column vector of A . We note the following two properties of A . One is that

$$(2) \quad c_{2i} + c_{2i+1} = c_1 \text{ for all } 1 \leq i \leq 5.$$

The other is that for $j_i \in J_i = \{2i, 2i + 1\}$ ($1 \leq i \leq 5$) we have

$$(3) \quad c_{j_1} + c_{j_2} + \dots + c_{j_5} = 0 \text{ iff } \sum j_i = \text{even}.$$

Let us check that \mathcal{F} is a $(11, \{0, 1, 2, 3, 5, 7\})$ -system. Choose $F, F' \in \mathcal{F}$ ($F \neq F'$) and set $I = \pi(F \cap F')$ where π is the projection. We shall show that $|I| \in \{0, 1, 2, 3, 5, 7\}$. By construction, if c_j is a linear combination of c_i , $i \in I$ then we have $c_j \in I$.

First suppose that $1 \in I$. Then by (2) we must have $|I \cap J_i| = 0$ or 2 for all $1 \leq i \leq 5$, in particular, $|I| = \text{odd}$. We can accept

$|I| \in \{1, 3, 5, 7\}$. Suppose that $|I| = 9$. Since $|I \cap J_i| \neq 1$, I contains precisely 4 of J_1, \dots, J_5 . But then by (3) and (2), I must contain all of J_1, \dots, J_5 , which is a contradiction.

Next we suppose that $1 \notin I$. Then by (2) we must have $|I \cap J_i| \leq 1$ for all $1 \leq i \leq 5$, in particular, $|I| \leq 5$. By (3) we cannot have $|I| = 4$. This concludes that \mathcal{F} is a $(11, \{0, 1, 2, 3, 5, 7\})$ -system.

By construction, we have $|\mathcal{F}| = (n/11)^5 = n^5/161051$ where $n = |V(\mathcal{F})| = 11 \cdot 2^d$. On the other hand, a $(11, \{0, 1, 2, 3, 5, 7\})$ -system with size $|\mathcal{F}| = n^5/(15 \cdot 2^{17}) = n^5/1966080$ was already constructed in [5]. (Unfortunately some of the values in the tables contained in [5] seem to be inaccurate.) Both constructions use

$$\mathcal{V} = \{(x, y, z, w, \lambda) \in \mathbb{F}_2^5 - \{0\} : x^2 + xy + y^2 + zw = 0\}.$$

In fact \mathcal{V} coincides with the set of column vectors of the matrix A .

6.2. The case $k = 12$. Let $L_0 = \{0, 1, 2, 3, 4, 6\}$. We consider (k, L) -systems with $k = 12$ and L listed below ($16 + 4 = 20$ cases).

- (I) $L_0, L_0 \cup \{7\}, L_0 \cup \{9\}, L_0 \cup \{10\}, L_0 \cup \{7, 9\}, L_0 \cup \{7, 10\}, L_0 \cup \{9, 10\}, L_0 \cup \{7, 9, 10\}, L_0 \cup \{12\}, L_0 \cup \{7, 12\}, L_0 \cup \{9, 12\}, L_0 \cup \{10, 12\}, L_0 \cup \{7, 9, 12\}, L_0 \cup \{7, 10, 12\}, L_0 \cup \{9, 10, 12\}, L_0 \cup \{7, 9, 10, 12\}$.
- (II) $L_0 \cup \{8\}, L_0 \cup \{8, 9\}, L_0 \cup \{8, 12\}, L_0 \cup \{8, 9, 12\}$.

By computer search, we found that $\text{rank}(12, L) = 6$ for all L in (I) and (II). As we saw in section 4 that $\alpha(12, L_0) = 6$, we now have $\alpha(12, L) = 6$ for all L in (I) and (II). The Steiner system $S(5, 6, 12)$ is the unique generator set of the corresponding intersection structure for all L in (I). There are precisely two intersection structures for all L in (II) — one is generated by $S(5, 6, 12)$ and the other is \mathcal{I}_{12} described below.

Let $J_i = \{2i - 1, 2i\}$ for $i = 1, \dots, 6$ and set

$$\mathcal{P} = \{J_a \cup J_b \cup J_c \cup J_d : \{a, b, c, d\} \in \binom{[6]}{4}\} \subset \binom{[12]}{8},$$

$$\mathcal{Q} = \{\{j_1, j_2, j_3, j_4, j_5, j_6\} : j_i \in J_i \text{ and } \sum j_i = \text{even}\} \subset \binom{[12]}{6}.$$

Then we define $\mathcal{I}_{12}^* = \mathcal{P} \cup \mathcal{Q}$ which is the generator set of \mathcal{I}_{12} . The corresponding generating matrix over $GF(2)$ is the following:

$$A = \begin{pmatrix} 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & 1 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 & 0 & 0 & 1 & 1 & 1 & 1 & 1 & 1 \\ 0 & 1 & 0 & 1 & 0 & 1 & 0 & 1 & 0 & 1 & 0 & 1 \end{pmatrix}.$$

Let \mathcal{F} be the 12-partite family generated by A with vertex partition $[n] = V_1 \cup \cdots \cup V_{12}$ where $V_i \cong \mathbb{F}_2^d$. Let c_i be the i -th column vector of A . We note the following two properties of A . One is that

$$(4) \quad (c_{2s-1} + c_{2s}) + (c_{2t-1} + c_{2t}) = 0 \text{ for all } 1 \leq s < t \leq 6.$$

The other is that for $j_i \in J_i = \{2i-1, 2i\}$ ($1 \leq i \leq 6$) we have

$$(5) \quad c_{j_1} + c_{j_2} + \cdots + c_{j_6} = 0 \text{ iff } \sum j_i = \text{even}.$$

Let us check that \mathcal{F} is a $(12, \{0, 1, 2, 3, 4, 6, 8\})$ -system. Choose $F, F' \in \mathcal{F}$ ($F \neq F'$) and set $I = \pi(F \cap F')$ where π is the projection. We shall show that $|I| \in \{0, 1, 2, 3, 4, 6, 8\}$. By construction, if c_j is a linear combination of c_i , $i \in I$ then we have $c_j \in I$.

If I contains one of J_1, \dots, J_6 , then by (4) we must have $|I \cap J_i| = 0$ or 2 for all i , in particular, $|I| = \text{even}$. We can accept $|I| \in \{0, 2, 4, 6, 8\}$. Suppose that $|I| = 10$. Then I contains precisely 5 of J_1, \dots, J_6 . But then by (5) and (4), I must contain all of J_1, \dots, J_6 , which is a contradiction.

Suppose now that $|I| = \text{odd}$. Then by (4) we need $|I \cup J_i| \leq 1$ for all $1 \leq i \leq 6$, which implies $|I| \leq 5$. But $|I| = 5$ is impossible because of (5).

Therefore, \mathcal{F} is a $(12, \{0, 1, 2, 3, 4, 6, 8\})$ -system with size $|\mathcal{F}| = (n/12)^6$ where $n = |V(\mathcal{F})| = 12 \cdot 2^d$.

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